

# Assignment 1

Thinking in Layers



#### Assignment Overview

- Getting familiar with Packets and Layers
- Doing some hands-on work with packets on the Internet
- Learning more about different application layer protocols



































#### Circuit Switching

- Propagation delay
- Transmission delay
- Queuing delay
- Nodal processing delay
- Setup delay

- Propagation delay
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- Queuing delay
- Nodal processing delay





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- The important key is circuit switching: The no sharing principle means that bandwidth must be reserved for a circuit/user,
- regardless of whether it is used!
- Hence: Some information we get is not relevant!

- Instead: 
$$|Users| = \frac{Available Bandwidth}{Bandwidth per User} \rightarrow |Users| = \frac{60mbit}{1mbit} = 60$$

# Question 1 (b)





#### - Hence: Some information we get is not relevant!

regardless of whether it is used!

- Instead:  $|Users| = \frac{Available Bandwidth}{Bandwidth per User} \rightarrow |Users| = \frac{60mbit}{1mbit} = 60$ 

- The important key is circuit switching: The no sharing principle

means that bandwidth must be reserved for a circuit/user,

- **Brief** justification: "Because of the no sharing principle, only the bandwidth users *may* need matters."







# Questions?



Data Networks

Assignment 1

























- OSI/ISO Layering Model
  - Layer 1:
  - Layer 2:
  - Layer 3:
  - Layer 4:
  - Layer 5:
  - Layer 6:
  - Layer 7:



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  - Layer 1:
  - Layer 2:
  - Layer 3:
  - Layer 4:
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  - Layer 7:





- OSI/ISO Layering Model
  - Layer 1: Physical Layer
  - Layer 2:
  - Layer 3:
  - Layer 4:
  - Layer 5:
  - Layer 6:
  - Layer 7:





- OSI/ISO Layering Model
  - Layer 1: Physical Layer
  - Layer 2: Link Layer
  - Layer 3:
  - Layer 4:
  - Layer 5:
  - Layer 6:
  - Layer 7:





- OSI/ISO Layering Model
  - Layer 1: Physical Layer
  - Layer 2: Link Layer
  - Layer 3: Network Layer
  - Layer 4:
  - Layer 5:
  - Layer 6:
  - Layer 7:





- OSI/ISO Layering Model
  - Layer 1: Physical Layer
  - Layer 2: Link Layer
  - Layer 3: Network Layer
  - Layer 4: Transport Layer
  - Layer 5:
  - Layer 6:
  - Layer 7:





- OSI/ISO Layering Model
  - Layer 1: Physical Layer
    - moves actual bits from one node to the next.
    - Works over the *physical* connection between two hosts
  - Layer 2: Link Layer
    - node-2-node (directly connected) data transfer.
    - Works independent of the underlying physical layer (WiFi, Copper, Optical)
  - Layer 3: Network Layer
    - Allows the exchange of data between hosts that are not directly connected.
  - Layer 4: Transport Layer
    - allows applications to exchange data
    - Enables transparent client-to-server data exchange.





# Questions?



Data Networks

Assignment 1

Figure 1 shows the three hosts H1, H2, and H3 that connect to the three switches S1, S2, and S3, respectively. All switches connect to the router R1. Explain the decapsulation performed by every device when H1 sends a packet to H3



Figure 1: Network topology


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Figure 1: Network topology





Data Networks





Data Networks





Data Networks





Data Networks





Data Networks







## Questions?



Data Networks

Select four universities in Germany, four universities in different European countries and four universities in different countries outside of Europe (12 in total). Ensure that they are pingable and record the average time it takes to get a reply from each target (i.e., the RTT).







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University Domain







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# UniversityDSaarland UniversityunWürzburgunTU MunichturTU Berlintu.

Domain uni-saarland.de uni-wuerzburg.de

tum.de

tu.berlin





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University	Domain
Saarland University	uni-saarland.de
Nürzburg	uni-wuerzburg.de
ՐՍ Munich	tum.de
۲U Berlin	tu.berlin
Dxford	ox.ac.uk
Porto	up.pt
Athens	uoa.gr
Stockholm	su.se





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#### RTT (ms)



- $> \sim$ \$ ping -c 4 uni-saarland.de
- > PING uni-saarland.de (134.96.7.179) 56(84) bytes of data.
- >64 bytes from webuni...land.de (134.96.7.179): icmp\_seq=1
  ttl=55 time=36.7 ms
- >64 bytes from webuni...land.de (134.96.7.179): icmp\_seq=2
  ttl=55 time=37.3 ms
- >64 bytes from webuni...land.de (134.96.7.179): icmp\_seq=3
   ttl=55 time=37.2 ms
- >64 bytes from webuni...land.de (134.96.7.179): icmp\_seq=4
  ttl=55 time=37.8 ms
- >--- uni-saarland.de ping statistics ---
- >4 packets transmitted, 4 received, 0% packet loss, time 3003ms
- >rtt min/avg/max/mdev = 36.718/37.277/37.830/0.393 ms



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Question 3 (a)					
Select four universities in Germany, four universities in different European countries and four universities in different countries outside of Europe (12 in total). Ensure that they are pingable and record the average time it takes to get a reply from each target (i.e., the RTT).	University Saarland University Würzburg TU Munich TU Berlin Oxford Oxford Porto Athens Stockholm MIT Sao Paulo Melbourne	Domain uni-saarland.de uni-wuerzburg.de tum.de tu.berlin ox.ac.uk up.pt uoa.gr su.se mit.edu usp.br unimelb.edu.au	RTT (ms) 37.277		
	Singapore	nus.edu.sg			



Data Networks

>ping -c 4 uni-wuerzburg.de

> PING wrz1114.rz.uni-wuerzburg.de (132.187.1.114) 56(84)
bytes of data.

>--- wrz1114.rz.uni-wuerzburg.de ping statistics ---

>4 packets transmitted, 0 received, 100% packet loss, time 3008ms







Data Networks



Germany, four

Select four universities in

University Domain Saarland University uni-saarland.de uni-luebeck.de tum.de tu.berlin ox.ac.uk up.pt uoa.gr su.se mit.edu usp.br unimelb.edu.au nus.edu.sg

RTT (ms) 37.277



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#### Question 3 (a)

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Lübeck	uni-luebeck.de	30.095
TU Munich	tum.de	
TU Berlin	tu.berlin	
Oxford	ox.ac.uk	
Porto	up.pt	
Athens	uoa.gr	
Stockholm	su.se	
MIT	mit.edu	
Sao Paulo	usp.br	
Melbourne	unimelb.edu.au	
Singapore	nus.edu.sg	



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	su.se	46.888
	mit.edu	16.251
	usp.br	235.717
	unimelb.edu.au	15.458
	nus.edu.sg	15.675



Determine your geographical distance (in kilometers) from the target. Can you find targets that have a smaller RTT but a larger distance than other targets? Why can this happen? (Give at least 2 reasons)

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Saarland University	uni-saarland.de	37.277		
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Data Networks

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Lübeck	uni-luebeck.de	30.095	575
TU Munich	tum.de	38.279	357
TU Berlin	tu.berlin	29.655	578
Oxford	ox.ac.uk	15.416	647
Porto	up.pt	61.082	1513
Athens	uoa.gr	52.542	1830
Stockholm	su.se	46.888	1329
MIT	mit.edu	16.251	5828
Sao Paulo	usp.br	235.717	9671
Melbourne	unimelb.edu.au	15.458	16450
Singapore	nus.edu.sg	15.675	10395

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Data Networks

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Data Networks

- –Technology: Faster Servers, Cold-Starts, Better Cables, Routing Quirks, …
- -Global distribution: The **responding** server can be geographically closer than the city (e.g. when using CDNs)
- Measurement: Average time distorted by outliers (e.g. random, queuing delays)




















### Question 3 (c)

- 1. import pandas as pd
- 2. import matplotlib.pyplot as plt
- 3. import seaborn as sns
- 4. df = pd.read\_csv('ping.csv')
- 5. fig, ax = plt.subplots()
- 6. sns.scatterplot(data=df, x='Distance'
  y='RTT', hue='University')
- 7. plt.show()







### Question 3 (c)

import pandas as pd

import matplotlib.pyplot as plt 2. Ping results of universities import seaborn as sns 3. University Saarland df = pd.read csv('ping.csv') 4. Lübeck 200 Munich Berlin fig, ax = plt.subplots() 5. Oxford sns.scatterplot(data=df, x='Distance' ິມ 150 6. Porto y='RTT', hue='University') Athens Average RTT Stockholm ax.set xlabel('Distance (km)') 7. MIT ax.set ylabel('Average RTT (ms)') 8. 100 Sao Paulo ax.set title('Ping results of 9. Melbourne universities') Singapore 10. plt.show() 50 7500 10000 0 2500 5000 12500 15000 Distance (km)



1.



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- 8. ax.set\_ylabel('Average RTT (ms)')
- 9. ax.set\_title('Ping results of universities')

- 12. plt.show()





Data Networks

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import pandas as pd

Question 3 (c)

- 4. df = pd.read\_csv('ping.csv')
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- 7. ax.set\_xlabel('Distance (km)')
- 8. ax.set\_ylabel('Average RTT (ms)')
- 9. ax.set\_title('Ping results of universities by region')

- 12. plt.show()





1.

2.

3.

Data Networks







Data Networks

## Question 4 (a)

• A sends a packet of 250 *bytes* to *C*, measuring an RTT of 20 ms. How many milliseconds of the RTT are made up of queueing delays?





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• A sends a packet of 250 *bytes* to *C*, measuring an RTT of 20 ms. How many milliseconds of the RTT are made up of queueing delays?



$$d_{total} = d_{proc} + 2d_{trans} + d_{queue} + d_{prop}$$





#### $\rightarrow 2 * d_{queue} = 10ms$ are made up of queuing delays



#### Data Networks

Assignment 1

## Question 4 (a)

- A sends a packet of 250 bytes to C, measuring an RTT of 20 ms.
- Assume a nodal processing delay of 0 ms
- Assume a bandwidth of  $2 \text{ Mbit/s} = 2 * 10^6 \text{ bit/s}$
- Link lengths of 1 km for  $L_1$ and 300 km for  $L_3$
- Assume a signal propagation speed of 10<sup>8</sup> m/s
- All replies have the same size as the outgoing packet

$$d_{total} = d_{proc} + 2d_{trans} + d_{queue} + d_{prop}$$

$$d_{queue} = d_{total} - d_{proc} - 2d_{trans} - d_{prop}$$

$$d_{total} = \frac{1}{2}RTT = \frac{1}{2} * 20ms = 10ms$$

$$d_{proc} = 0ms$$

$$d_{trans} = \frac{Size}{Bandwidth} = \frac{8 * 250 \ bit}{2 * 10^6 \ \frac{bit}{s}} = 1ms$$

$$d_{prop} = \frac{L_1 + L_3}{Speed} = \frac{301km}{10^8 m/s} \approx 3ms$$

$$d_{queue} = 10ms - 0ms - 2ms - 3ms = 5ms$$







Data Networks

## Question 4 (b)

- *C* sends a packet to *A*, measuring an RTT of 65 *ms*.
- What was the approximate size of the packet assuming empty queues this time?

$$d_{trans} = \frac{Size}{Bandwidth}$$

$$d_{total} = d_{proc} + 2d_{trans} + d_{queue} + d_{prop}$$

$$d_{trans} = \frac{(d_{total} - d_{proc} - d_{queue} - d_{prop})}{2}$$



## Question 4 (b)

- *C* sends a packet to *A*, measuring an RTT of 65 *ms*.
- What was the approximate size of the packet assuming empty queues this time?
- Assume a nodal processing delay of 0 *ms*
- The propagation delay stays unchanged 3 ms
- Assume a bandwidth of  $2 \text{ Mbit/s} = 2 * 10^6 \text{ bit/s}$

$$d_{trans} = \frac{Size}{Bandwidth} = \frac{Size}{2 * 10^{6} bit/s}$$

$$d_{total} = d_{proc} + 2d_{trans} + d_{queue} + d_{prop}$$

$$d_{trans} = \frac{\left(d_{total} - d_{proc} - d_{queue} - d_{prop}\right)}{2}$$

$$d_{trans} = \frac{\left(32.5ms - 0ms - 0ms - 3ms\right)}{2}$$

$$d_{trans} = 14.75ms = \frac{2}{size}$$

$$d_{trans} = 29500 bit \approx 3,687 byte$$

 $\rightarrow$  The packet size is 3687 byte





## Question 4 (c)

- *B* sends a packet of 1000 bytes to *C* and measures an *RTT* of 32 *ms*.
- Calculate a theoretical upper bound for the length of the line  $L_2$ .

$$\begin{split} d_{prop} &= \frac{L_2 + L_3}{Speed} \\ d_{total} &= d_{proc} + 2d_{trans} + d_{queue} + d_{prop} \\ d_{prop} &= d_{total} - d_{proc} - 2d_{trans} - d_{queue} \end{split}$$





## Question 4 (c)

- *B* sends a packet of 1000 bytes to *C* and measures an *RTT* of 32 *ms*.
- Calculate a theoretical upper bound for the length of the line  $L_2$ .
- Assume a nodal processing delay of 0 ms
- Assume a queueing delay of 0 ms
- Assume a signal propagation speed of 10<sup>8</sup> m/s
- Link lengths of 1 km for L<sub>1</sub> and 300 km for L<sub>3</sub>

$$\begin{aligned} d_{prop} &= \frac{L_2 + L_3}{Speed} \\ d_{total} &= d_{proc} + 2d_{trans} + d_{queue} + d_{prop} \\ d_{prop} &= d_{total} - d_{proc} - 2d_{trans} - d_{queue} \\ d_{prop} &= 16ms - 0ms - \frac{8 * 1000 \text{ bit}}{2 * 10^6 \text{ bit/s}} - 0ms \end{aligned}$$

$$d_{prop} = 8ms$$

 $L_2 = 500 km$ 

 $\rightarrow$  L<sub>2</sub> can be up to 500 km long







Data Networks

## Question 5 (a)



Choose an application layer protocol that is NOT listed below: SMTP, POP3, IMAP, SSH, FTP(S), HTTP(S), DNS.

Start some research on the chosen protocol and try to briefly discuss it along the following lines:

a) In 3-4 sentences, summarize its purpose and basic functionality

NTP is a protocol that allows the synchronization of system clocks in a network.

NTP Servers sync to a reference clock or other servers, NTP clients sync their internal clocks to servers.

NTP also specifies a format to exchange messages containing the times and algorithms on how to adjust the internal clock.



## Question 5 (b)



Choose an application layer protocol that is NOT listed below: SMTP, POP3, IMAP, SSH, FTP(S), HTTP(S), DNS.

Start some research on the chosen protocol and try to briefly discuss

it along the following lines:

b) Which transport layer protocol does it use? Is the application layer protocol standardized? Can you find the standard?

NTP uses UDP NTPv4 (RFC5905) is a proposed standard NTPv3 (RFC1305) is a draft standard







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## Feedback?



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