

# Homework 9

Multimedia & VoIP, Datacenter, SDN



#### Homework Overview

- Multimedia Applications & VoIP
- Datacenter Networks
- Software Defined Network (SDN)





## **Question 1: Multimedia Applications**



Please provide a short answer for the questions below.







Multimedia applications can be classified into three categories. Name and briefly describe each category in 2-3 sentences. State how much they have tolerance to delay, jitter, and packet loss.





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#### **Streaming Stored Audio and Video**





**Streaming Stored Audio and Video** 

- prerecorded video/audio
- placed on the server
- requested on demand by the users





**Streaming Stored Audio and Video** 

- prerecorded video/audio
- placed on the server
- requested on demand by the users
- not sensitive to Delay / jitter / packet loss





#### **Conversational Voice- and Video-over-IP**





**Conversational Voice- and Video-over-IP** 

- •real-time
- similar to the traditional circuit-switched telephone service





**Conversational Voice- and Video-over-IP** 

- real-time
- similar to the traditional circuit-switched telephone service
- sensitive to delay and jitter
- occasional loss only causes occasional glitches in audio/video playback





#### Streaming Live Audio and Video









Streaming Live Audio and Video

- similar to traditional broadcast radio/television
- a live radio/television transmission to the users







**Streaming Live Audio and Video** 

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- a live radio/television transmission to the users
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What is the client playout delay? Explain how it can affect the users' experience in two ways.

#### Answers:

A certain amount of time which client buffered the received data till it renders and plays it.





What is the client playout delay? Explain how it can affect the users' experience in two ways.

Answers:

Too large: users experience a large delay at any new start





What is the client playout delay? Explain how it can affect the users' experience in two ways.

Answers:

**Too large:** users experience a large delay at any new start **Too small:** users might experience multiple refreshing (lags) during rendering.







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Answers:

 Reliable delivery of frames is of little use if retransmissions take too long to keep the conversation going





What is the motivation behind using FEC instead of transport/application layer reliability in VoIP traffic?

Answers:

- Reliable delivery of frames is of little use if retransmissions take too long to keep the conversation going
- FEC-like mechanisms which **provide recovery without retransmission** appear to be better-suited





## Questions?



Data Networks

Multimedia & VoIP, Datacenter, SDN

Question 2(a): Video Streaming

Why are video frames categorized into three types (I-Frames, P-Frames, and B-Frames)?





Question 2(a): Video Streaming



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Question 2(a): Video Streaming

Why are video frames categorized into three types (I-Frames, P-Frames, and B-Frames)?

#### Answer:

Video is a flow of frames, in order to compress the data of a video we can divide frames into three categories as I,P and B frame.





# What is the difference between these three types of frames?







# What is the difference between these three types of frames?



Question 2(b): Video Streaming



What is the difference between these three types of frames?

Answers:

I frames are the most important ones that do not rely on any other frame they are sent periodically.



Question 2(b): Video Streaming



What is the difference between these three types of frames?

Answers:

I frames are the most important ones that do not rely on any other frame they are sent periodically.

**P frames** are predicted frames. These frames contain the changes from proceeding I-frames.



Question 2(b): Video Streaming



What is the difference between these three types of frames?

Answers:

I frames are the most important ones that do not rely on any other frame they are sent periodically.

**P frames** are predicted frames. These frames contain the changes from proceeding I-frames.

**B frames** are bidirectional frame, these frames are relative to the past and future I/P frame.





## Questions?



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## Question 3: Hands-on lab for VoIP

You need Wireshark1 traffic analyzer and a real traffic trace (.pcap files) in order to solve following questions. Do not use any other tools.

In this assignment we set up a lab with a PBX2 and two clients. Clients (Bob and Alice) are able to use VoIP service and PBX is the service provider.

Wireshark provides a set of analysis tools for VoIP traffic, most of them are in the statistics and telephony menus in the main window of the Wireshark. Familiarize yourself with these tools, you will need them to solve all of the following questions.

Important note: Your answers must include description for your reasoning and how did you reach to the answers. In cases you use display filters for a question, provide a snapshot of your filter in your answer sheet.

The VoIP lab topology is as Figure 1. We have traces from two vantage points. One is on Bob side which is the Bob.pcap3 file and the other one is the PBX.pcap3.



## Question 3: Hands-on lab for VoIP

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The VoIP lab topology is as Figure 1. We have traces from two vantage points. One is on Bob side which is the Bob.pcap3 file and the other one is the PBX.pcap3.



# Question 3: Hands-on lab for VoIP Alice Bob /OIP 'OIP PBX

#### Figure 1: VoIP lab topology.



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Based on trace and wireshark analysis fill Table 1 and put a summary on how you identified which IP belongs to which node: (use PBX.pcap file)

	PBX	Alice	Bob
IP Address			

Table 1: IP Addresses







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	PBX	Alice	Bob
IP Address			

Table 1: IP Addresses





No.	Time	Source	Destination	Protocol	Length  Info
	1 0.000000	192.168.0.236	192.168.0.133	SIP	563 Request: REGISTER sip:192.168.0.133 (1 binding)
	2 0.000302	192.168.0.133	192.168.0.236	SIP	617 Status: 401 Unauthorized
	3 0.000568	192.168.0.236	192.168.0.133	SIP	724 Request: REGISTER sip:192.168.0.133 (1 binding)
	4 0.000843	192.168.0.133	192.168.0.236	SIP	639 Status: 200 OK (REGISTER) (1 binding)
	5 2.800046	192.168.0.18	192.168.0.133	SIP	529 Request: REGISTER sip:192.168.0.133 (1 binding)
	6 2.800227	192.168.0.133	192.168.0.18	SIP	558 Status: 401 Unauthorized
	7 2.819991	192.168.0.18	192.168.0.133	SIP	693 Request: REGISTER sip:192.168.0.133 (1 binding)

Internet Protocol Version 4, Src: 192.168.0.236, Dst: 192.168.0.133
User Datagram Protocol, Src Port: 51028, Dst Port: 5060
Session Initiation Protocol (REGISTER)
Request-Line: REGISTER sip:192.168.0.133 SIP/2.0
~ Message Header

#### For the 1st packet:

	Message Header
	Via: SIP/2.0/UDP 192.168.0.236:51028; rport; branch=z9hG4bKPjdf56f9bccf1b
	Max-Forward <u>s: 70</u>
icket:	From: <sip:alice@192.168.0.133>;tag=141594470f0048689de668a682f67042</sip:alice@192.168.0.133>
	> To: <sip:alice@192.168.0.133></sip:alice@192.168.0.133>
	Call-ID: 0ad934a3d32d452bbb70b412bee5f040
	[Generated Call-ID: 0ad934a3d32d452bbb70b412bee5f040]
	> CSeq: 35068 REGISTER
	User-Agent: MicroSIP/3.20.5
	<pre>&gt; Contact: <sip:alice@192.168.0.236:51028:ob></sip:alice@192.168.0.236:51028:ob></pre>





No.	Time	Source	Destination	Protocol	Length  Info
	1 0.000000	192.168.0.236	192.168.0.133	SIP	563 Request: REGISTER sip:192.168.0.133 (1 binding)
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#### For t

	~ Message Header				
	Via: SIP/2.0/UDP 192.168.0.236:51028; rport; branch=z9hG4bKPjdf56f9bccf1k				
	Max-Forwards: 70				
che 1st packet:	<pre>&gt; From: <sip:alice@192.168.0.133>;tag=141594470f0048689de668a682f67042</sip:alice@192.168.0.133></pre>				
1	<pre>&gt; To: <sip:alice@192.168.0.133></sip:alice@192.168.0.133></pre>				
	Call-ID: 0ad934a3d32d452bbb70b412bee5f040				
	[Generated Call-ID: 0ad934a3d32d452bbb70b412bee5f040]				
	CSeq: 35068 REGISTER				
	User-Agent: MicroSIP/3.20.5				
	<pre>Contact: <sip:alice@192.168.0.236:51028:ob></sip:alice@192.168.0.236:51028:ob></pre>				



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No.	Time	Source	Destination	Protocol	Length  Info
	1 0.000000	192.168.0.236	192.168.0.133	SIP	563 Request: REGISTER sip:192.168.0.133 (1 binding)
	2 0.000302	192.168.0.133	192.168.0.236	SIP	617 Status: 401 Unauthorized
	3 0.000568	192.168.0.236	192.168.0.133	SIP	724 Request: REGISTER sip:192.168.0.133 (1 binding)
	4 0.000843	192.168.0.133	192.168.0.236	SIP	639 Status: 200 OK (REGISTER) (1 binding)
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	7 2.819991	192.168.0.18	192.168.0.133	SIP	693 Request: REGISTER sip:192.168.0.133 (1 binding)

	<ul> <li>Internet Protocol Version 4, Src: 192.168.0.18, Dst: 192.168.0.133</li> <li>User Datagram Protocol, Src Port: 5060, Dst Port: 5060</li> <li>Session Initiation Protocol (REGISTER)</li> </ul>
	Request-Line: REGISTER sip:192.168.0.133 SIP/2.0
	✓ Message Header
	> Via: SIP/2.0/UDP 192.168.0.18;branch=z9hG4bK232cc8c1-82cca759
	<pre>→ From: sip:bob@192.168.0.133;tag=f961d7f7-10a453ee</pre>
For the 5th packet:	> To: sip:bob@192.168.0.133
for the junpacket.	Call-ID: 8a75dc04-d11a3a9e-a6c70db1-820eaeb4
	[Generated Call-ID: 8a75dc04-d11a3a9e-a6c70db1-820eaeb4]
	> CSeq: 278 REGISTER
	<pre>&gt; Contact: <sip:bob@192.168.0.18;x-reg=e6c9f60392c33cd9>;expires=600</sip:bob@192.168.0.18;x-reg=e6c9f60392c33cd9></pre>
	Content-Length: 0
	Expires: 600
	User-Agent: Sipnetic/1.0.37 Android





No.	Time	Source	Destination	Protocol	Length  Info
	10.000000	192.168.0.236	192.168.0.133	SIP	563 Request: REGISTER sip:192.168.0.133 (1 binding)
	2 0.000302	192.168.0.133	192.168.0.236	SIP	617 Status: 401 Unauthorized
	3 0.000568	192.168.0.236	192.168.0.133	SIP	724 Request: REGISTER sip:192.168.0.133 (1 binding)
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Answers:

	PBX	Alice	Bob
IP Address	192.168.0.133	192.168.0.236	192.168.0.18

IP Addresses







# Why do Alice and Bob send register requests twice? What is the difference between them? (use PBX.pcap file)







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Answers:

The first register is replied by 401 Unauthorized from the server side, this packet contains WWW-Authenticate header.





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Answers:

The first register is replied by 401 Unauthorized from the server side, this packet contains WWW-Authenticate header.

By this header, in the second register, by this header the user can decide about encryption method of its communications password.





Why do Alice and Bob send register requests twice? What is the difference between them? (use PBX.pcap file)

To make the following graph: Statistics - >Flow Graph

192.168.0.236









Call Reference	Caller (Name - Number 700x)	Calee (Name - Number 700x)	Who hangs up the call	Duration (seconds)

Table 2: Calls list







Call Reference	Caller (Name - Number 700x)	Calee (Name - Number 700x)	Who hangs up the call	Duration (seconds)

Table 2: Calls list







Apply the filter "sip.Method == "BYE"

sip.Me	p.Method == "BYE"								
<b>)</b> .	Time	Source	Destination	Protocol	Length  Info				
	210 68.336003	192.168.0.18	192.168.0.133	SIP	551 Request: BYE sip <mark>:7001</mark> @192.168.0.133:5060				
	215 68.337916	192.168.0.133	192.168.0.236	SIP	500 Request: BYE sip: <u>alic</u> e@192.168.0.236:51028;ob				
	646 185.498786	192.168.0.236	192.168.0.133	SIP	420 Request: BYE sip: <mark>7002</mark> @192.168.0.133:5060				
	652 185.520100	192.168.0.133	192.168.0.18	SIP	509 Request: BYE sip:bob@192.168.0.18;x-reg=E6C9F60392C33CD9				
	1115 264.641059	192.168.0.18	192.168.0.133	SIP	551 Request: BYE sip: <mark>7001</mark> @192.168.0.133:5060				
	1120 264.643473	192.168.0.133	192.168.0.236	SIP	500 Request: BYE sip:alice@192.168.0.236:50366;ob				







Telephony -> VoIP Calls

Start Time \land	Stop Time	Initial Speaker	From	То	Protocol	Duration	Packets	State	Comments
27.250257	68.336352	192.168.0.18	si <mark>r</mark> :bob@ <mark>192.168.0.133</mark>	sip:7001@192.168.0.133	SIP	00:00:41	12	COMPLETED	INVITE 401 200
27.293800	68.338184	192.168.0.133	"bob" <sip:7002@192.168.0.133></sip:7002@192.168.0.133>	<sip:alice@192.168.0.236:51028;ob></sip:alice@192.168.0.236:51028;ob>	SIP	00:00:41	13	COMPLETED	INVITE 200 200 200
141.472914	185.499258	192.168.0.236	<sip alice@192.168.0.133=""></sip>	<sip:7002@192.168.0.133></sip:7002@192.168.0.133>	SIP	00:00:44	12	COMPLETED	INVITE 401 200
141.477552	185.538902	192.168.0.133	"alice" <sip:7001@192.168.0.133></sip:7001@192.168.0.133>	<sip:bob@192.168.0.18;x-reg=e6c9f60392c33cd9></sip:bob@192.168.0.18;x-reg=e6c9f60392c33cd9>	SIP	00:00:44	12	COMPLETED	INVITE 200 200 200
223.073683	264.641276	192.168.0.18	sip bob@192.168.0.133	sip:7001@192.168.0.133	SIP	00:00:41	12	COMPLETED	INVITE 401 200
223.118813	264.643698	192.168.0.133	"bob" <sip:7002@192.168.0.133></sip:7002@192.168.0.133>	<sip:alice@192.168.0.236:50366;ob></sip:alice@192.168.0.236:50366;ob>	SIP	00:00:41	13	COMPLETED	INVITE 200 200 200







Answers:

Call Reference	Caller (Name - Number 700x)	Calee	Who hangs the call	Duration (s)
1	Bob 7002	Alice 7001	Bob	41
2	Alice 7001	Bob 7002	Alice	44
3	Bob 7002	Alice 7001	Bob	41

#### Calls list





### Questions?



Data Networks

Multimedia & VoIP, Datacenter, SDN



A new Datacenter topology design, Jellyfish, is to construct a random graph as its topology at the top-of-rack (ToR) switch layer, as Figure 2a. Each ToR switch has some number k of ports, of which it uses r to connect to other ToR switches, and uses the remaining k-r ports for servers. In this question, we assume there are two data centers with the same number of equipment (switches and servers). One uses Fat tree topology and the other one uses Jellyfish topology.





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#### **Question 4: Datacenter**



Figure 2: Datacenter topologies



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Fat-tree















# What are the bisection widths of two data centers given in Figure 2?

#### Fat tree bisection width: 8



















Jellyfish







## What are the bisection widths of two data centers given in Figure 2?

#### Jellyfish bisection width: 6



Jellyfish



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Answers:

• Jellyfish costs more.







Answers:

- Jellyfish costs more.
- Jellyfish is more complex and harder to maintain and rewire.







Answers:

- Jellyfish costs more.
- Jellyfish is more complex and harder to maintain and rewire.
- Jellyfish has higher cable cost. It needs longer cable / more options of cables length because switches are not necessarily connected to switches nearby.



### Question 4 (c)



In the random graph, there are usually more than one link between two nodes. But equal-cost multipath routing (ECMP) protocol doesn't work well in the Jellyfish datacenter. Why? What other problems may occur if we simply deploy current routing algorithms to the Jellyfish datacenter?



### Question 4 (c)



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### Question 4 (c)

Answers:

• ECMP doesn't work well because multiple paths between two nodes in a random graph (or Jellyfish) usually don't have equal costs.




## Question 4 (c)

Answers:

• ECMP doesn't work well because multiple paths between two nodes in a random graph (or Jellyfish) usually don't have equal costs.

• Most traffic between two nodes would go with the shortest path which leads to traffic unbalance.





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Paths with different costs (hops) may cause packet reordering.

- Most traffic between two nodes would go with the shortest path which leads to traffic unbalance.
- ECMP doesn't work well because multiple paths between two nodes in a random graph (or Jellyfish) usually don't have equal costs.

## Question 4 (c)

Answers:





## Questions?



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• Decouple software from hardware





- Decouple software from hardware
- Decouple control plane from data plane





- Decouple software from hardware
- Decouple control plane from data plane
- Faster protocol/feature updating





- Decouple software from hardware
- Decouple control plane from data plane
- Faster protocol/feature updating
- Provides an open standard API or HW/SW interface





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#### What are the functions of data plane and control plane?







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What are the functions of data plane and control plane? Answers:

• Data plane: packet streaming (e.g. forward, filter, buffer, mark, rate-limit, and measure packets).







What are the functions of data plane and control plane? Answers:

- Data plane: packet streaming (e.g. forward, filter, buffer, mark, rate-limit, and measure packets).
- Control plane: track topology changes, compute routes, install forwarding rules



### Question 5 (c): SDN



In 2049, we ran out of IPv6 addresses because robots learned how to replicate themselves. However, a genius invented IPv8, for which an address has 256 bits and the processing algorithm takes a new field, "INET", as input. If you were a network architect in the future world, which SDN technologies we learned in Data Network lectures would you choose to implement IPv8 protocol? Why would you choose it over the other?



### Question 5 (c): SDN



In 2049, we ran out of IPv6 addresses because robots learned how to replicate themselves. However, a genius invented IPv8, for which an address has 256 bits and the processing algorithm takes a new field, "INET", as input. If you were a network architect in the future world, which **SDN technologies** we learned in Data Network lectures would you choose to implement IPv8 protocol? Why would you choose it over the other?





• I will choose P4.





## Question 5 (c): SDN

- I will choose P4.
- IPv8 address is 256-bit which is different from the current IP address. With P4, programmers can **define new formats** while OpenFlow only supports existing formats.



## Question 5 (c): SDN

- I will choose P4.
- IPv8 address is 256-bit which is different from the current IP address. With P4, programmers can **define new formats** while OpenFlow only supports existing formats.
- IPv8 needs a new processing algorithm than the current IP processing algorithm. With P4, programmers can define new processing algorithms but OpenFlow only supports existing actions and there are only limited number of them.





## Questions?



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